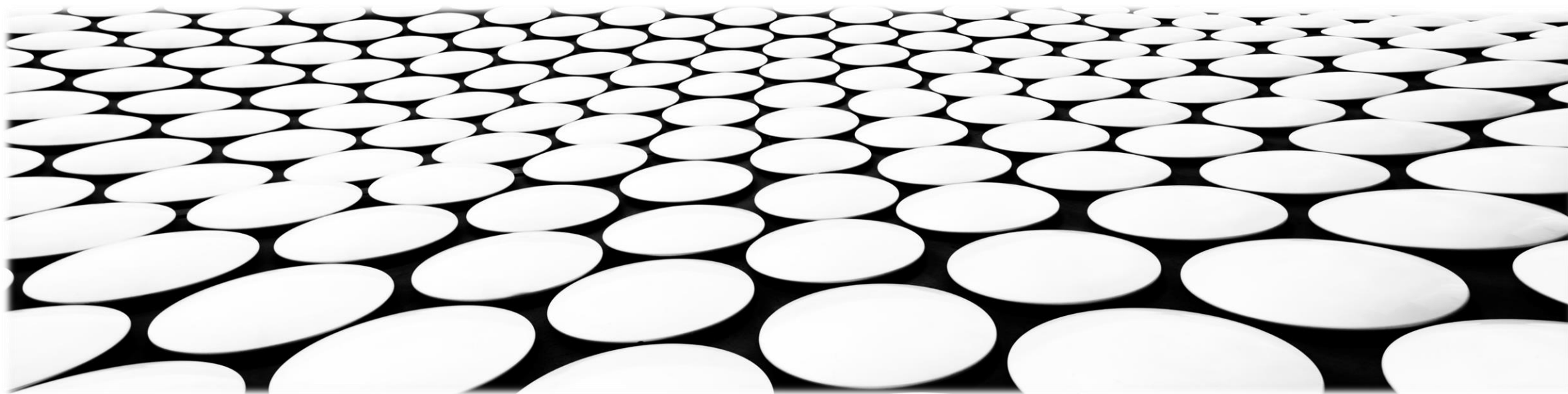

FUZZING TVM RELAY

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WHY DON'T WE HAVE MORE TVM TESTS?

	Lines of code (KSLOC)
Implementation (tvm/src, tvm/python)	233
Tests (tvm/tests)	141

- Test cases are program fragments
 - Tedious to write by hand
 - Complex interactions between features
 - Shapes need to match up
- Fuzzing could help

RELAY FUZZING APPROACH

- How do we generate Relay programs we know are valid?
- Use typing information: Given a type, generates expression fulfilling it
- We have a prototype! Only ~2000 lines of Python
- Supports most statically typed Relay constructs, ~20 operators

BASIC CASES IN FUZZING RELAY

- Most of Relay's type system plays nice
- Set goal type and work backwards
 - All types have a literal for a base case*
 - Connectives (let bindings, etc.) combine existing terms
- Ensuring termination: Fall back to a literal!

*This can get tricky with arbitrary ADTs.

Type-Product

$$\frac{\forall i \in [1, n]: \Delta; \Gamma \vdash p_i : T_i}{\Delta; \Gamma \vdash (p_1, \dots, p_n) : (T_1, \dots, T_n)}$$

Type-Projection

$$\frac{\Delta; \Gamma \vdash p : (T_1, \dots, T_n) \quad i \in [0, n]}{\Delta; \Gamma \vdash p.i : T_{i+1}}$$

Type-Let

$$\frac{\Delta; \Gamma \vdash v : T \quad \Delta; \Gamma, id : T \vdash e : T'}{\Delta; \Gamma \vdash \text{let } \%id = v; e : T'}$$

Type-Ref

$$\frac{\Delta; \Gamma \vdash n : T}{\Delta; \Gamma \vdash \text{Ref } n : \text{RefType}[T]}$$

Type-Read-Ref

$$\frac{\Delta; \Gamma \vdash r : \text{RefType}[T]}{\Delta; \Gamma \vdash !r : T}$$

Type-Write-Ref

$$\frac{\Delta; \Gamma \vdash r : \text{RefType}[T] \quad \Delta; \Gamma \vdash v : T}{\Delta; \Gamma \vdash r := v : ()}$$

THE TOUGH PART: SOLVING TYPE RELATIONS

- Type system includes constraints on tensor shapes!
- Argument types (shapes) affect result type (shape)
- Every single op has a relation!
- Hardest part: Implemented *imperatively* in C++

DEALING WITH TYPE RELATIONS: SOLVER-BASED APPROACH

- Encode type relations in a solver domain (e.g., ILP)
- Given return type, use solver to generate valid argument types
- Pro: Only one solver query at a time, easily composable
- Cons:
 - The solver is a dependency
 - Need to formalize the type relations in the solver domain

DEALING WITH TYPE RELATIONS: STOCHASTIC APPROACH

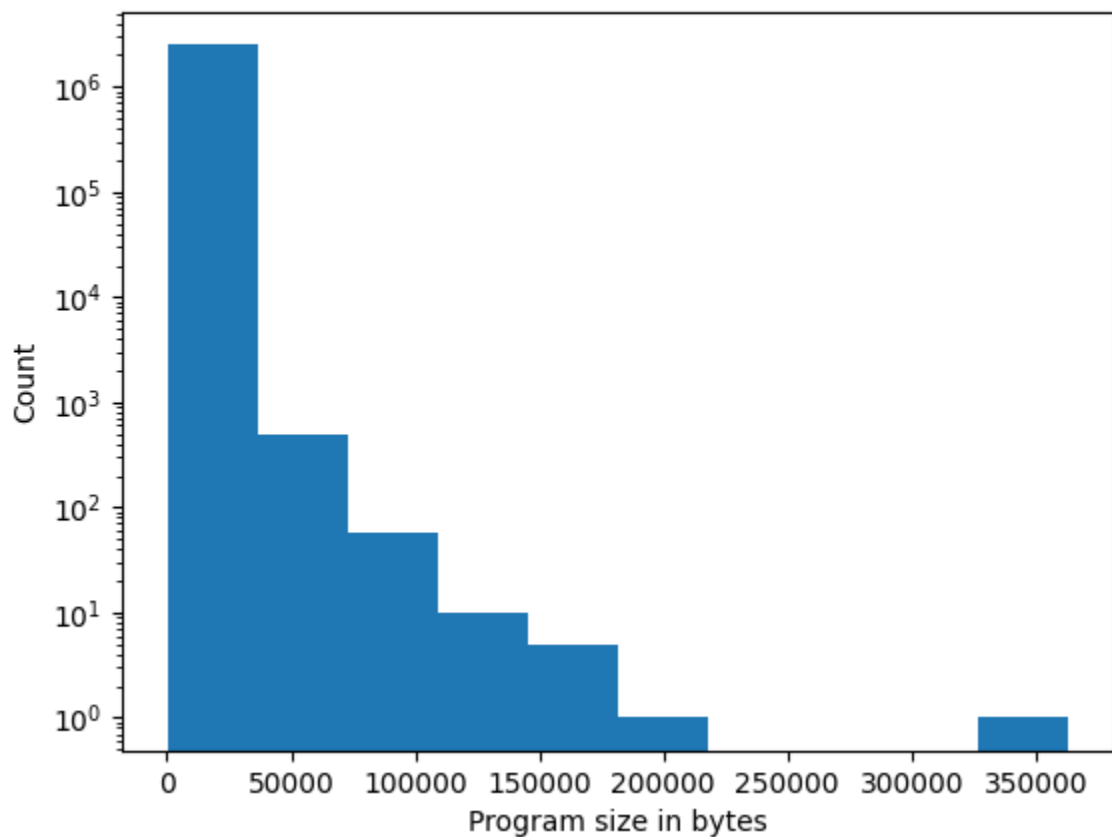
- Sample possible inputs, check which solutions work, keep a cache
- Use argument type–return type pairs to guide type generation
- Pro: Can reuse existing type relation implementations, no solver
- Con: Not as flexible as solver-based approach

BUGS FOUND

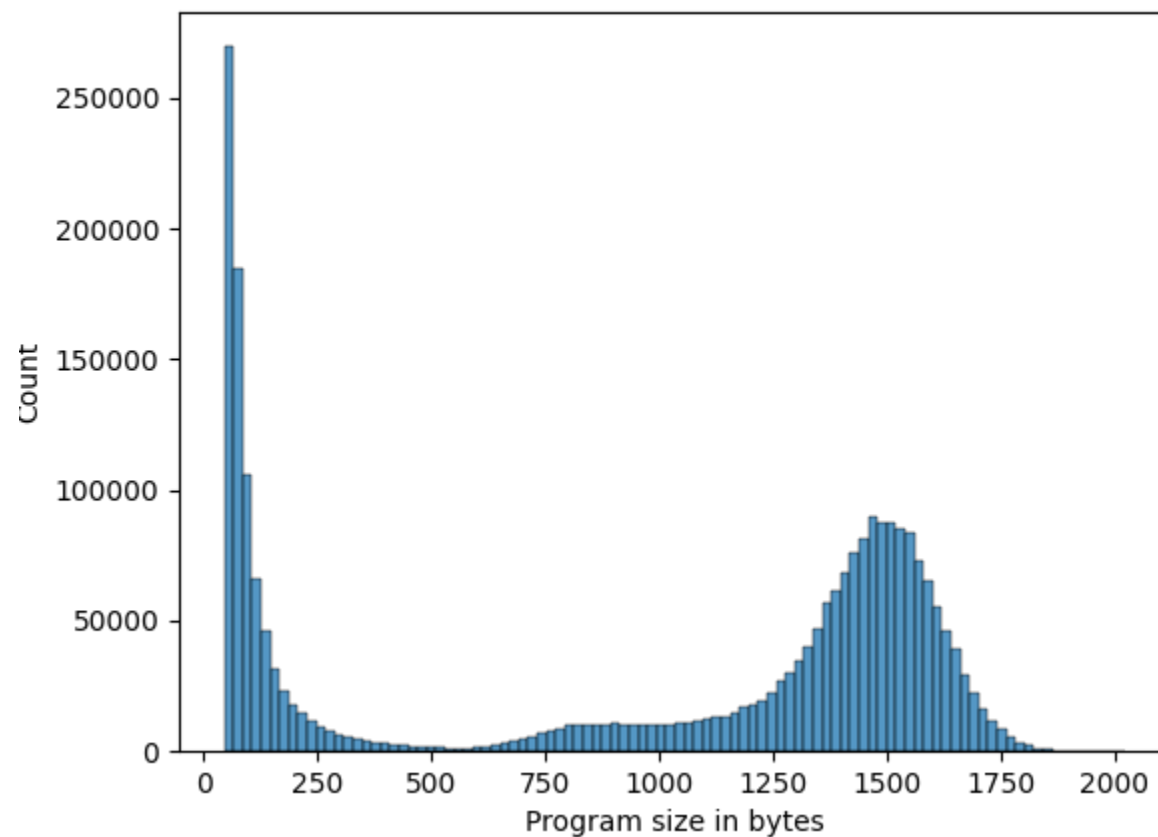
- Match exhaustion bug:
 - Found by fuzzer very quickly in small-scale test runs
 - Fix merged <https://github.com/apache/tvm/pull/7459>
- Missing bounds check in bias add specification:
 - Found manually while formalizing the type relation
 - Fix merged <https://github.com/apache/tvm/pull/7554>
- Also found a bug parsing refs of refs (fix not yet PR'd)

GENERATED PROGRAM SIZES

Type-Directed

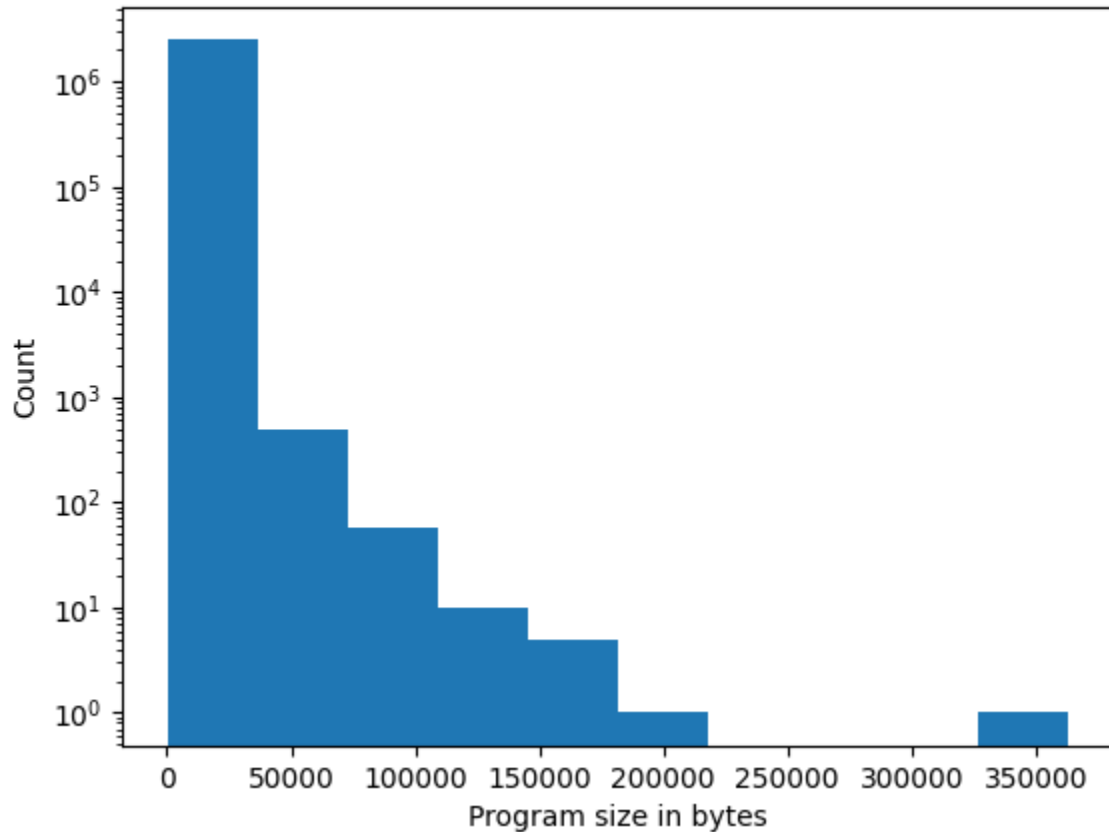


Grammar-Based

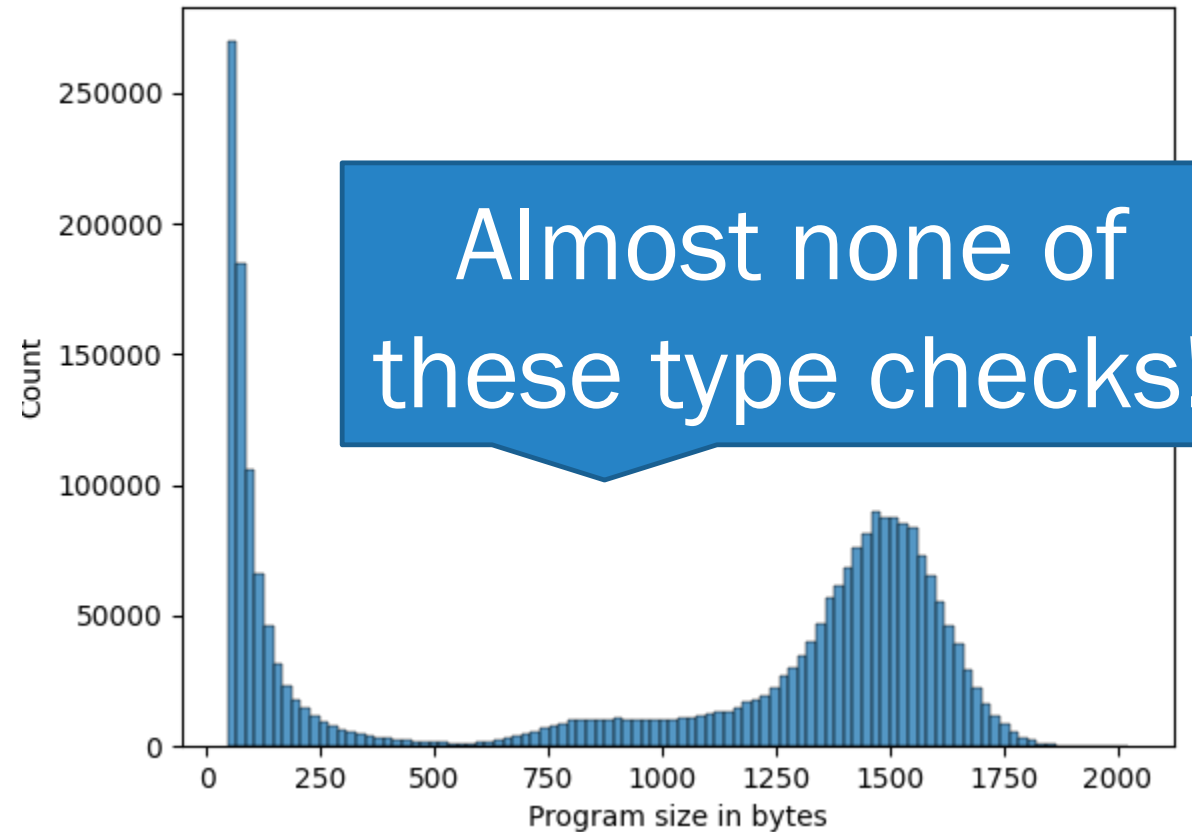


GENERATED PROGRAM SIZES

Type-Directed



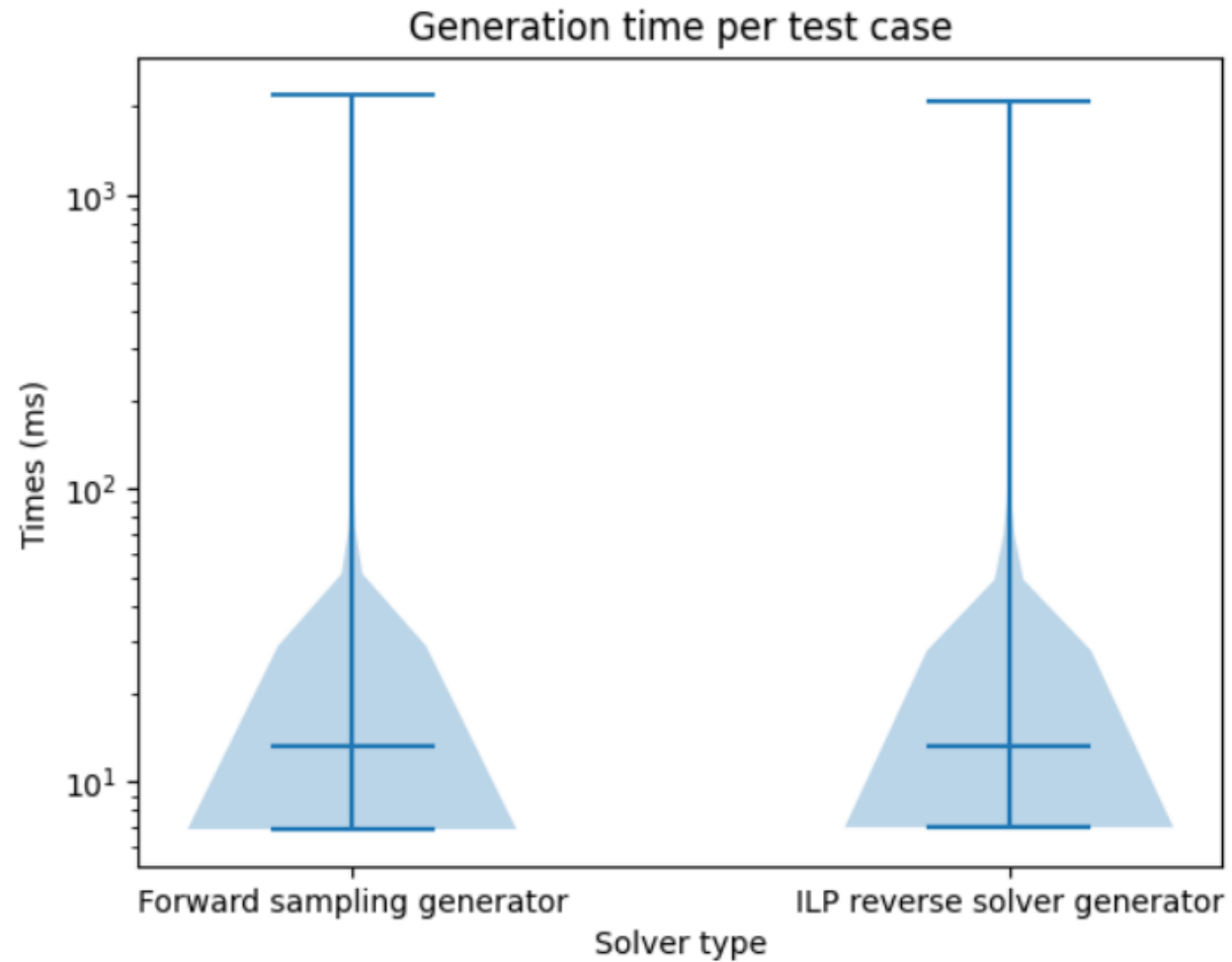
Grammar-Based



(SOME OF THE TRIVIAL CASES THAT DID TYPE CHECK)

```
def @main() -> () { match? 75 {} }
def @main() -> () { () }
def @main() -> () { match? (((),)) {} }
def @main() -> () { %793 = 8 ; () }
def @main() -> uint16 { match? () {} }
```

GENERATION SPEED



THE FUTURE OF THE FUZZER

- Prototype available at: https://github.com/slyubomirsky/relay_fuzzer
- Will create a TVM RFC for discussing the future of fuzzing
- Questions for the future:
 - How can we support dynamic or parametric shapes?
 - What testing oracles make the most sense to use?
 - How should we express constraints on generated programs?
 - What about mutating and minimizing Relay programs for bug reports?



THANK YOU!

